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Federated Bandit: A Gossiping Approach

Zhaowei Zhu*[†], Jingxuan Zhu*[§], Ji Liu[§], Yang Liu[†]

[†]UC Santa Cruz, [§]Stony Brook University

{zwzhu,yangliu}@ucsc.edu,{jingxuan.zhu,ji.liu}@stonybrook.edu

(*Equal contributions)



Motivating Example

- Scenario: Multiple hospitals test different treatment plans (arms)

Treatment effectiveness (arm 1)

Hospital 1: + ■ - □

Hospital 2: + ■ - □

Observations

■ ■ ■ ■ □ □ □

■ ■ □ □ □ □ □

Rewards

4/7

2/7

True (global) reward:
3/7

- Problem: Biased data → Biased reward → Biased decision
- Main focus: Data heterogeneity

Federated Bandit

In Multi-Armed Bandit (MAB) settings, we consider:

➤ Heterogeneity

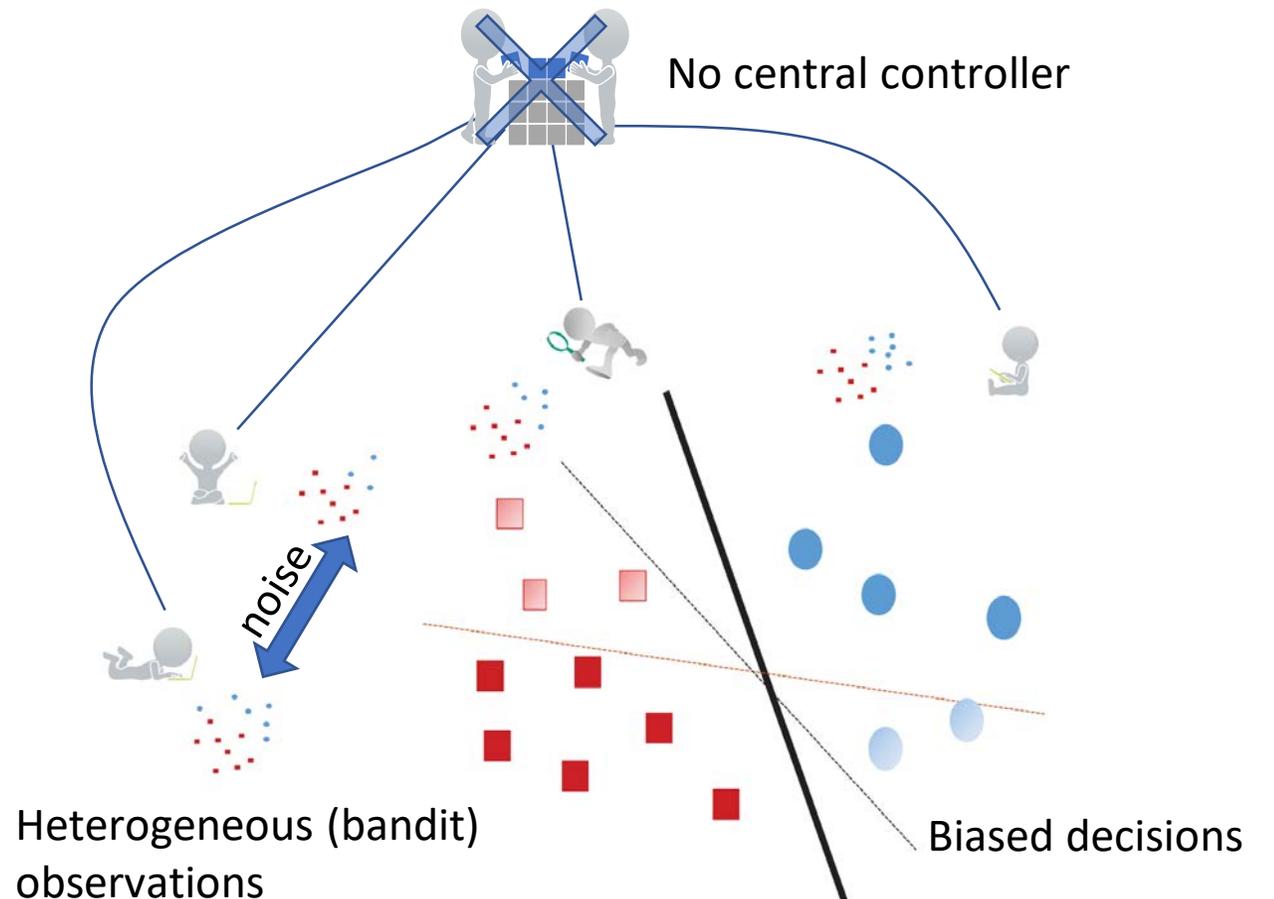
- Bias in local learning
- Problem may not be solved

➤ Decentralization

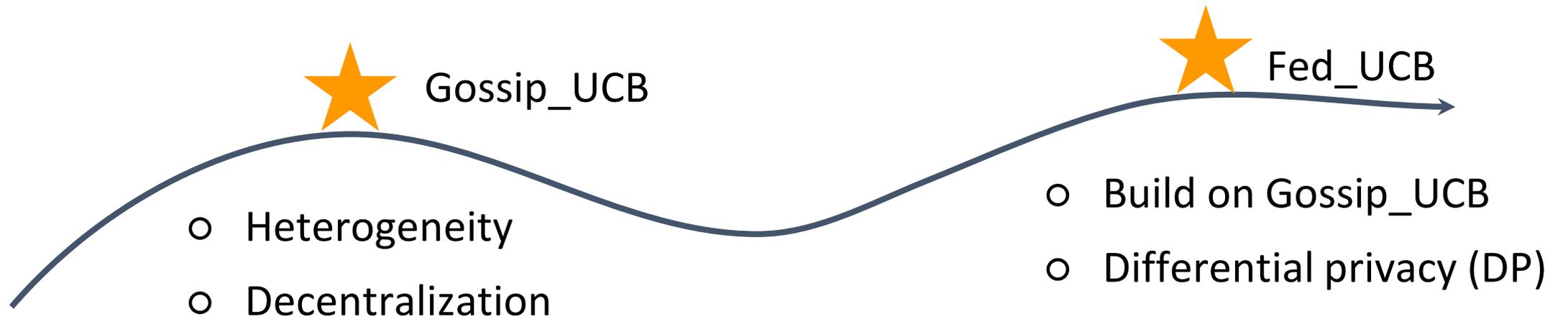
- No central-controller

➤ Privacy

- Protect agents' privacy in the worst cases during federation

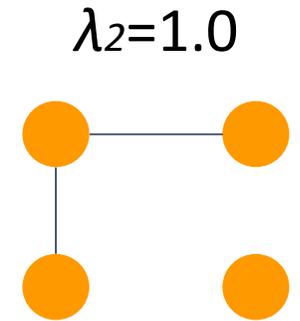
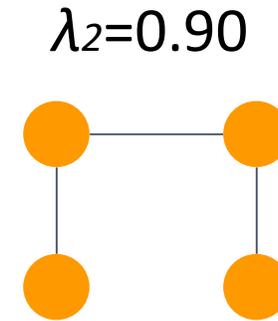
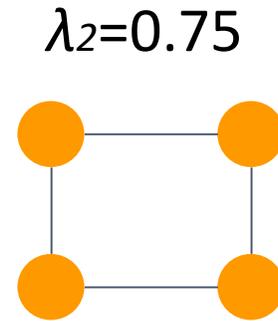
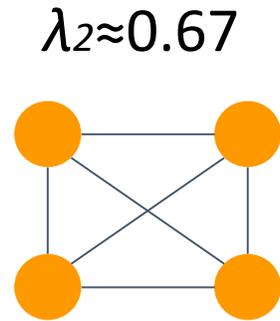


Road Map



Regret

λ_2 : 2nd-largest eigenvalue of the gossip matrix



➤ Gossip_UCB ($\propto \log T$, $1/\text{connectivity}$)

$$O(\max\{\text{poly}(N, M) \log T, \text{poly}(N, M) \log_{\lambda_2^{-1}} N\})$$

➤ Fed_UCB (ϵ -DP, $\propto \log T$, $1/\text{connectivity}$, $1/\epsilon$)

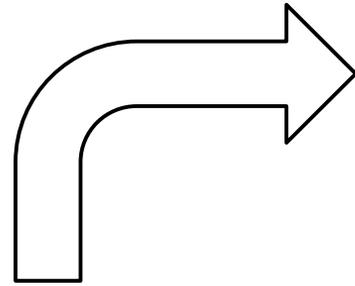
$$O(\max\{\frac{\text{poly}(N, M)}{\epsilon} \log^{2.5} T, \text{poly}(N, M) (\log_{\lambda_2^{-1}} N + \log T)\})$$

N : # agents

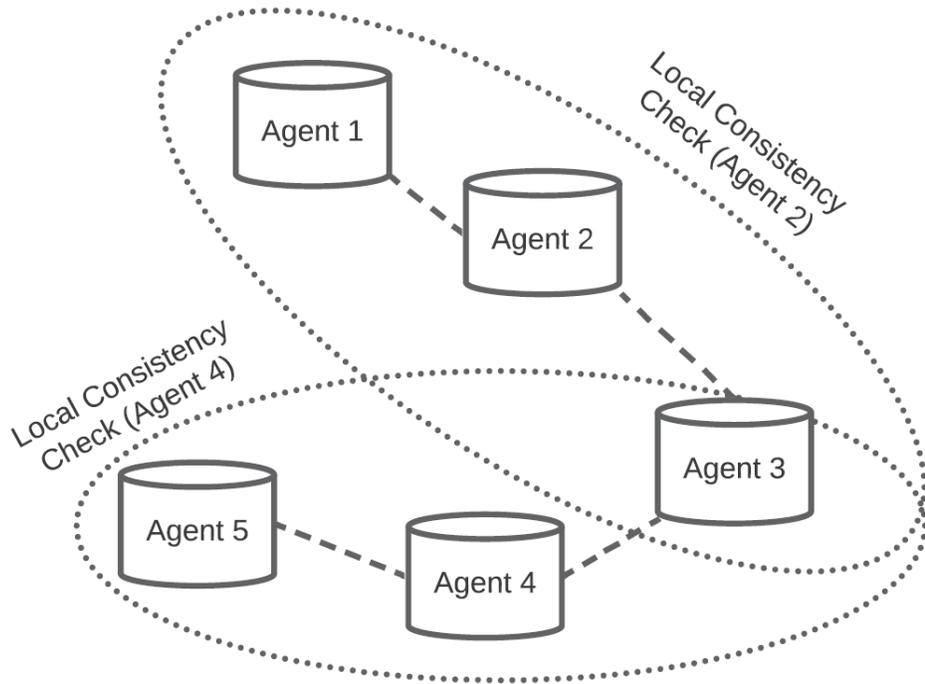
M : # arms

T : total time

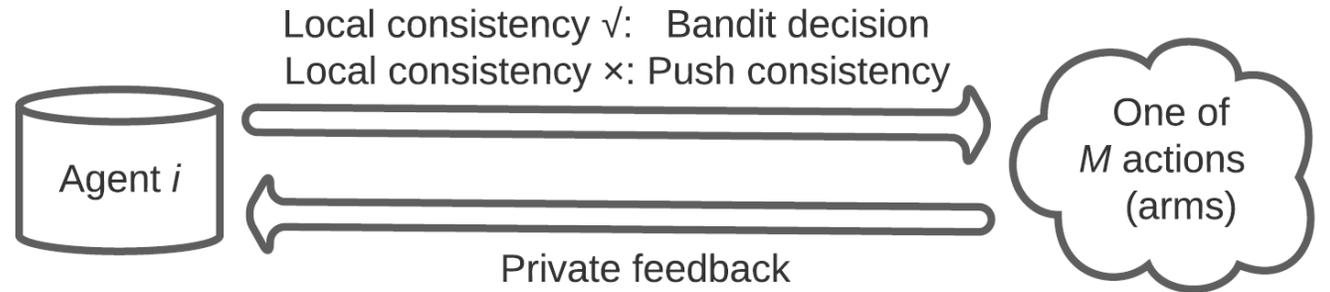
Algorithm



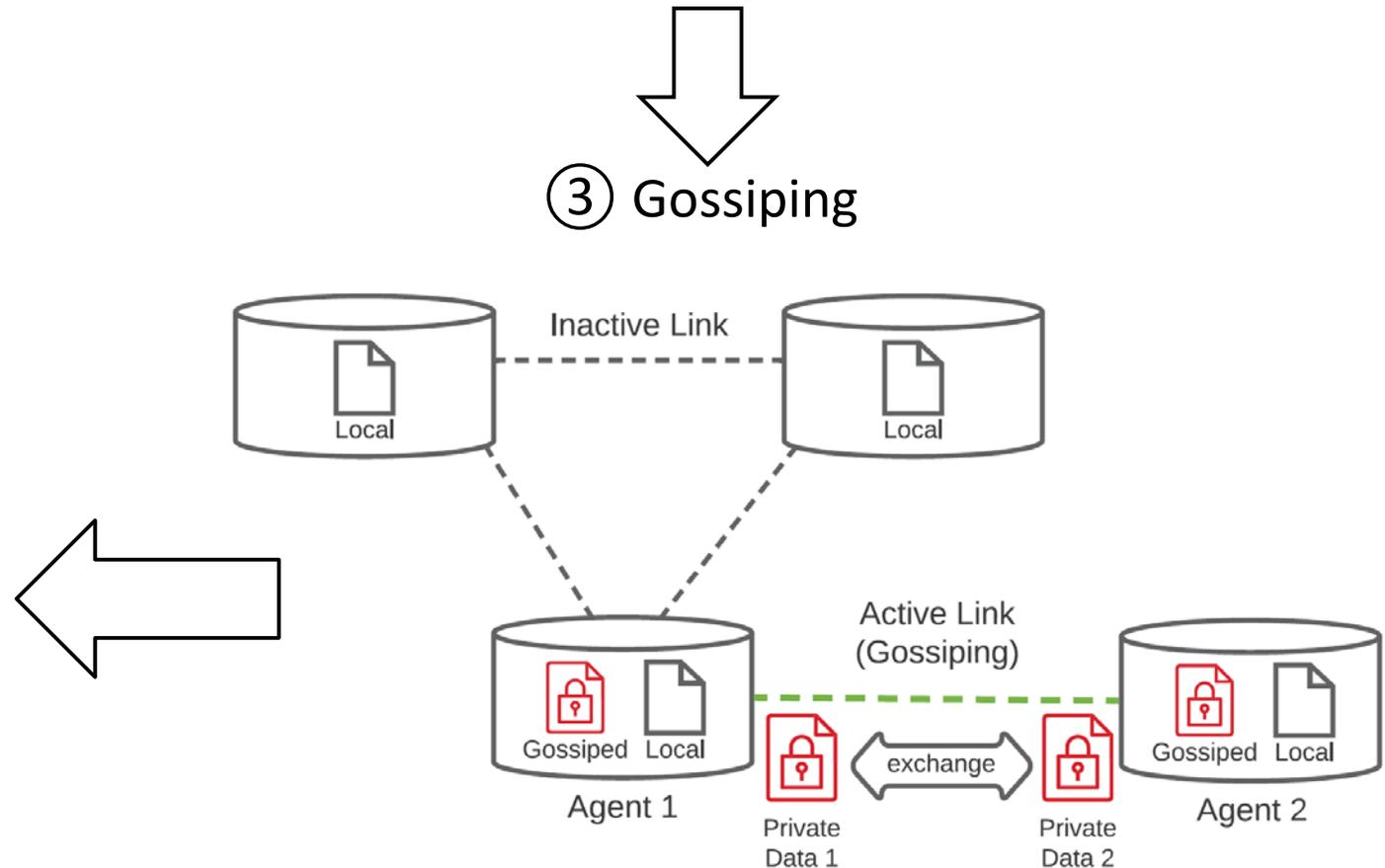
① Local consistency check



② Locally consistent decision making



③ Gossiping



Gossip_UCB



Gossip_UCB

- Heterogeneity
- Decentralization



Fed_UCB

- Build on Gossip_UCB
- Differential privacy (DP)

Gossip_UCB: Bandit Problem

➤ MAB with N agents, M arms (total time T):

- Agent i pulls arm $a_i(t)$ at time t , gets feedback $X_{i,a_i(t)}(t)$
- Ideal (global) feedback $X_{a_i(t)}(t)$
- Expectation of feedbacks: Local mean $\mu_{i,k}$
- Expectation of ideal feedbacks: Global mean μ_k
- Regret (μ_1 is the max):

$$R_i(T) = T\mu_1 - \sum_{t=1}^T \mathbb{E} [X_{a_i(t)}(t)]$$

➤ Homogenous setting:

- Agent gets exact feedback $\mu_{i,k} = \mu_k$

➤ Heterogenous setting:

- Agent gets biased/noisy feedback $\mu_{i,k} \neq \mu_k$

Gossip_UCB: Heterogeneous Feedbacks

- **Global** mean μ_k vs. **local** mean $\mu_{i,k}$:

$$\mu_k := \frac{1}{N} \sum_{i=1}^N \mu_{i,k}$$

- Estimates:

- **Sample mean (MEAN)** $\tilde{X}_{i,k}(t)$: estimate of **local** mean (averaged observations)
- **Estimate of rewards (EST)** $\vartheta_{i,k}(t)$: estimate of **global** mean

- Action: $a_i(t) = \arg \max_k \vartheta_{i,k}(t-1) + \boxed{C_{i,k}(t)}$ → **Upper confidence bound (UCB)**

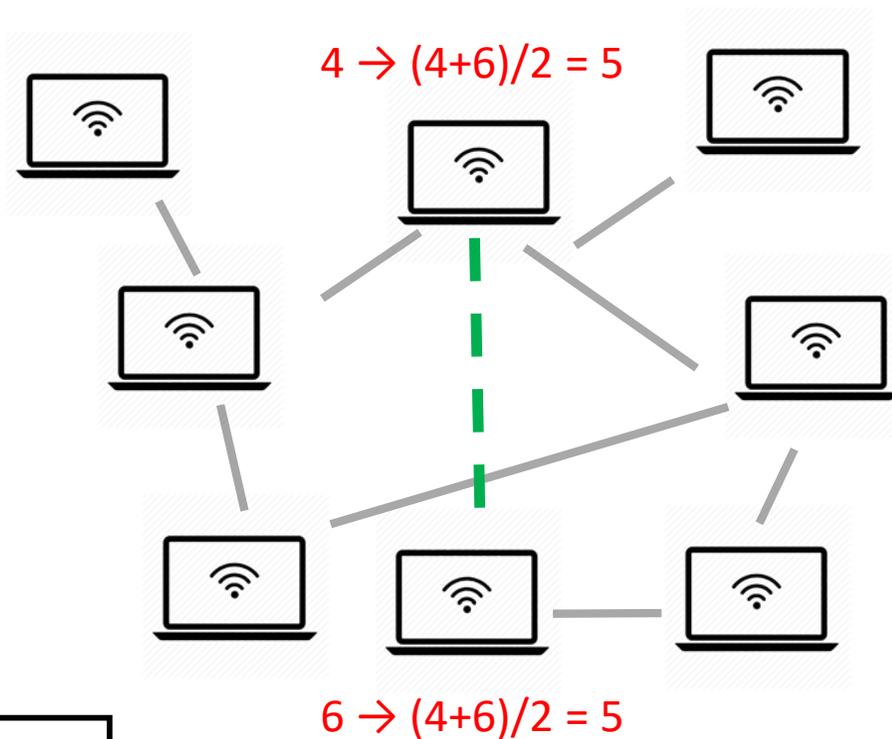
- Heterogeneity:

- **Local** mean \neq **Global** mean \rightarrow **Sample mean** \neq **Estimate of rewards**

Design it!

Gossip_UCB: Gossiping

Communication among agents (**classical gossiping** [1]):



Active Link — —
Idle Link —

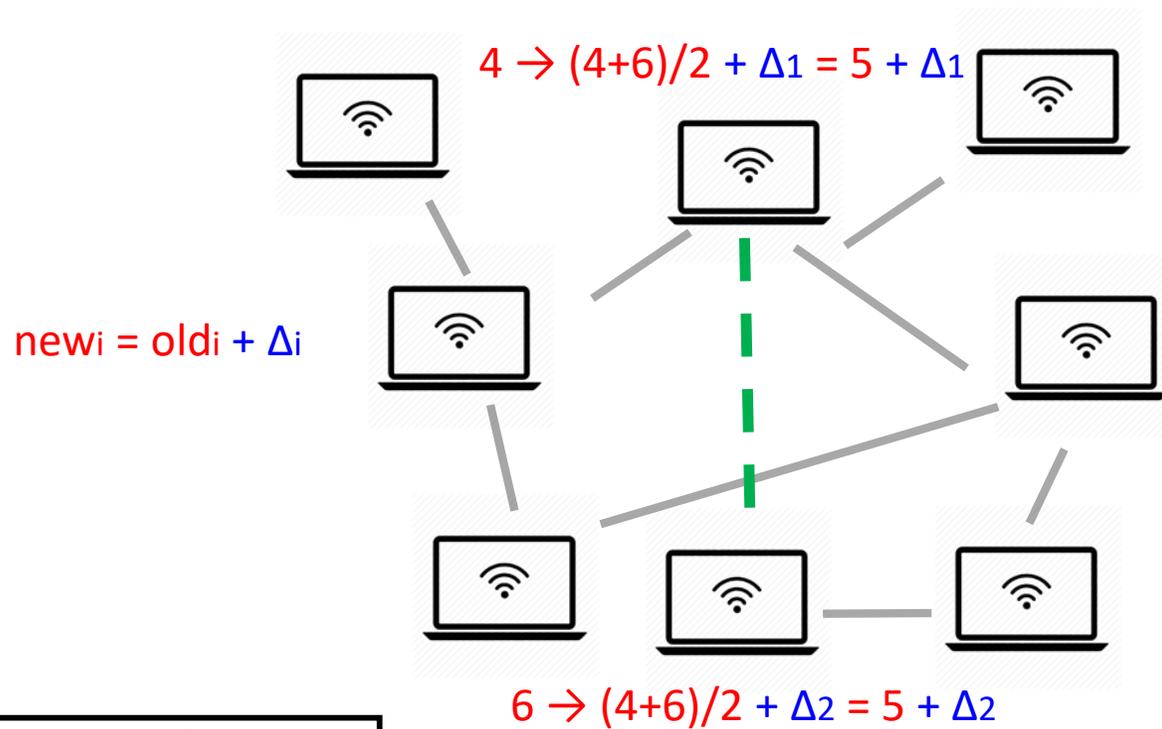
- Features:
- One edge activated at each time t
 - Selected agents on the edge **exchange information**
 - Others **do not update**

$$\text{GOSSIP} = (\text{EST}_{\text{self}} + \text{EST}_{\text{other}})/2$$

[1] Stephen Boyd, Arpita Ghosh, Balaji Prabhakar, and Devavrat Shah. 2006. Randomized Gossip Algorithms. *IEEE TIT* 52, 6 (2006), 2508–2530.

Gossip_UCB: Gossiping

Communication among agents (**bandit gossiping** [2]):



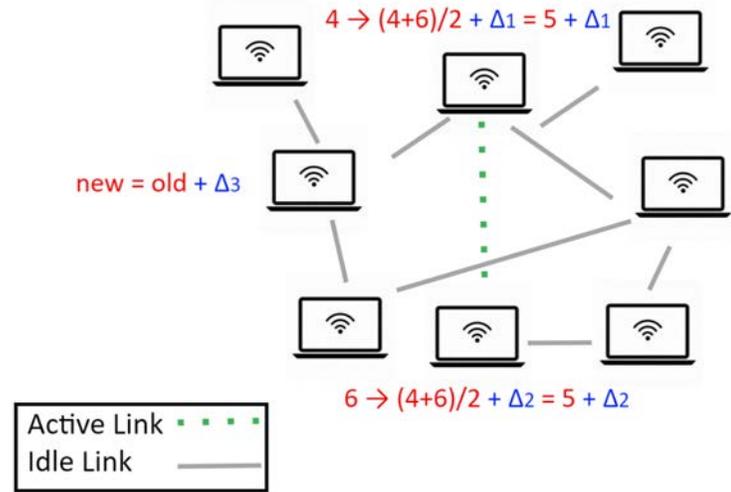
- Features:
- One edge activated at each t
 - Selected agents on the edge **exchange information + gradient**
 - Others: **gradient update**

Active Link - - -
Idle Link —

[2] Yang Liu, Ji Liu, and Tamer Başar. 2018. Differentially private gossip gradient descent. In *IEEE CDC*. IEEE, 2777–2782.

Gossip_UCB: Gossiping

Communication among agents (**bandit gossiping**):



Each agent:

$$EST_t = \begin{cases} \text{GOSSIP} + MEAN_t - MEAN_{t-1} & \text{(Gossiping update)} \\ EST_{t-1} + MEAN_t - MEAN_{t-1} & \text{(Normal update)} \end{cases}$$

if agent i is selected to gossip with agent j **then**

agent i sends $\vartheta_{i,k}(t-1)$ to agent j

agent i receives $\vartheta_{j,k}(t-1)$ from agent j

$$\vartheta_{i,k}(t) := \frac{\vartheta_{i,k}(t-1) + \vartheta_{j,k}(t-1)}{2} + \tilde{X}_{i,k}(t) - \tilde{X}_{i,k}(t-1)$$

else

$$\vartheta_{i,k}(t) := \vartheta_{i,k}(t-1) + \tilde{X}_{i,k}(t) - \tilde{X}_{i,k}(t-1)$$

end

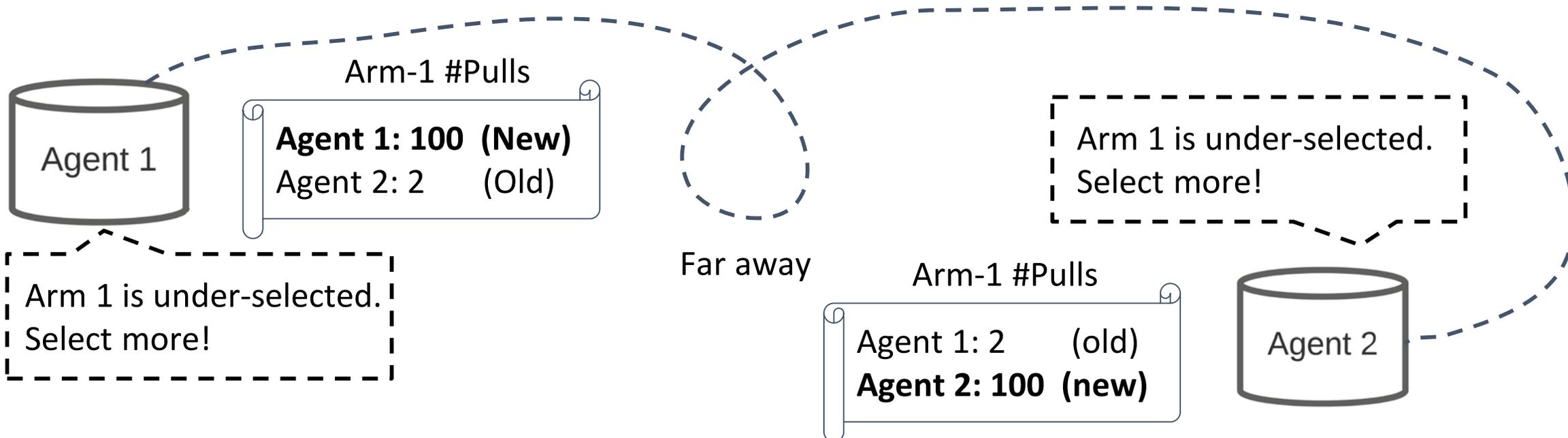
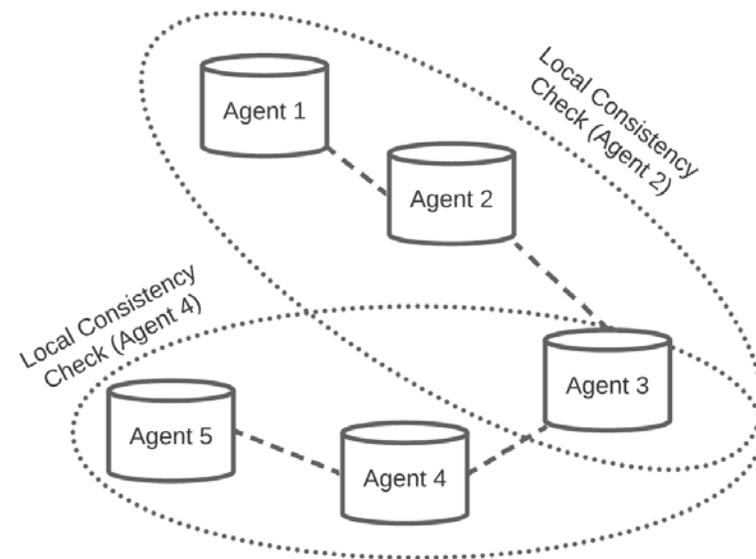
// gossiping update

// normal update

Gossip_UCB: Consistency

➤ Problem:

- Arm true estimates depend on global estimates (average of local estimates)
- #Pull affects the calculation of UCB



Gossip_UCB: Consistency

➤ Local consistency:

- #LocalPulls is close to the local estimate of #GlobalPulls (Lemma 2)
- Local estimate is not a bottleneck

➤ Global consistency:

- Max #LocalPulls $\leq 2 \times$ #LocalPulls_i, for all agent i (Lemma 3)

➤ Summary:

- Bound the information inconsistency due to propagation delay
- Facilitate a fully-decentralized solution

Gossip_UCB: Concentration Bound

➤ Concentration Bound for Local Estimates (Theorem 1)

- With some conditions and a high probability $(1 - \bar{p}_0)$:

where

$$\mathbb{P}(|\vartheta_{i,k}(t) - \mu_k| \geq C_{i,k}(t)) < \frac{2}{t^2}$$

$$C_{i,k}(t) = \sqrt{\frac{2N}{n_{i,k}(t)} \log t + \alpha_1}$$

- Notations: **local estimates**, **global mean**, **#LocalPulls** (global consistency)
- Challenges:
 - Coupling effects of gossiping and bandit learning

Gossip_UCB: Proof Overview

- Guarantees on the Consistency among Agents
 - Information propagation (Lemma 1)
 - Actual local consistency (Lemma 2)
 - Global consistency (Lemma 3)
- Concentration Bound for Local Estimates
 - Bound $|LocalEst - E[LocalEst]|$ (Lemma 4)
 - Bound $|E[LocalEst] - GlobalMean|$ (Lemma 5)
- Regret Upper Bound for Gossip_UCB (Theorem 1)

Fed_UCB



Gossip_UCB

- Heterogeneity
- Decentralization



Fed_UCB

- Build on Gossip_UCB
- Differential privacy (DP)

Fed_UCB: Differential Privacy

➤ Why necessary?

- Directly leaking some information that might appear to be “anonymized” can be used to cross-reference with other datasets to breach privacy [3]
- Worst-case privacy guarantee

➤ Differential privacy (DP) [4]:

A (randomized) algorithm \mathcal{B} is ϵ -differentially private if for any adjacent streams $\{X_{i,k}(t)\}_{t=1}^T$ and $\{X'_{i,k}(t)\}_{t=1}^T$, and for all sets $\mathcal{O} \in \mathcal{C}$,

$$\mathbb{P} \left[\mathcal{B}(\{X_{i,k}(t)\}_{t=1}^T) \in \mathcal{O} \right] \leq e^\epsilon \cdot \mathbb{P} \left[\mathcal{B}(\{X'_{i,k}(t)\}_{t=1}^T) \in \mathcal{O} \right].$$

[3] Latanya Sweeney. 2000. Simple demographics often identify people uniquely. *Health (San Francisco)* 671, 2000, 1–34.

[4] Cynthia Dwork, Frank McSherry, Kobbi Nissim, and Adam Smith. 2006. Calibrating noise to sensitivity in private data analysis. *In Theory of cryptography conference*. Springer, 265–284.

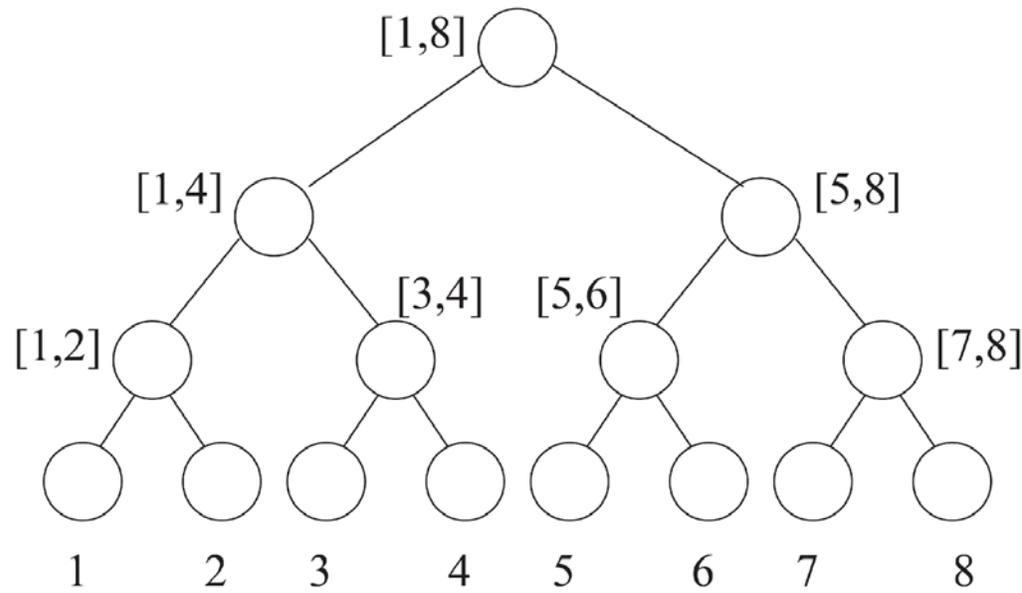
Fed_UCB: Online DP

➤ Naive method:

- Add Laplacian noise $\text{Lap}(T/\epsilon)$ to each observation

➤ Partial sum [5]:

- Add Laplacian noise $\text{Lap}((\log T)/\epsilon)$ following a tree structure



Example:

Node 4: $\text{Noise}_{[1,4]}$

Node 6: $\text{Noise}_{[1,4]} + \text{Noise}_{[5,6]}$

Node 7: $\text{Noise}_{[1,4]} + \text{Noise}_{[5,6]} + \text{Noise}_7$

Fed_UCB: Concentration Bound

➤ Concentration Bound for Local Estimates

- Guarantee ϵ -DP, with some conditions and a high probability $(1 - \frac{2N}{n_{i,k}(t)} - p_0)$:

$$\mathbb{P}(|\tilde{\vartheta}_{i,k}(t) - \mu_k| \geq \tilde{C}_{i,k}(t)) < \frac{2}{t^2}$$

where

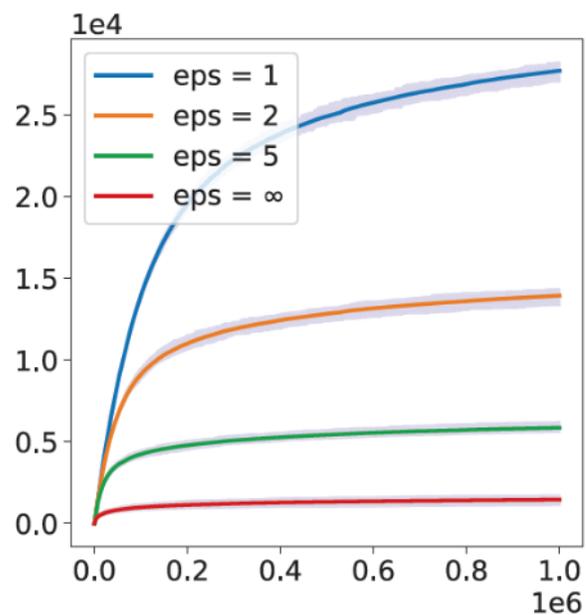
$$\tilde{C}_{i,k}(t) = \alpha_1 + \sqrt{2N \left(\frac{128N \log^2 T \cdot \log t \cdot \log n_{i,k}(t)}{n_{i,k}^2(t) \epsilon^2} + \frac{1}{n_{i,k}(t)} \right) \log t.}$$

➤ Compared with Gossip_UCB: two changes

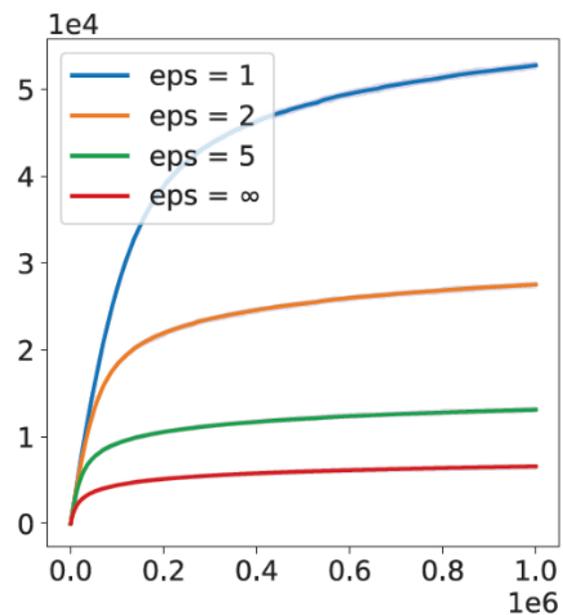
- Upper confidence bound
- Probability

Experiments

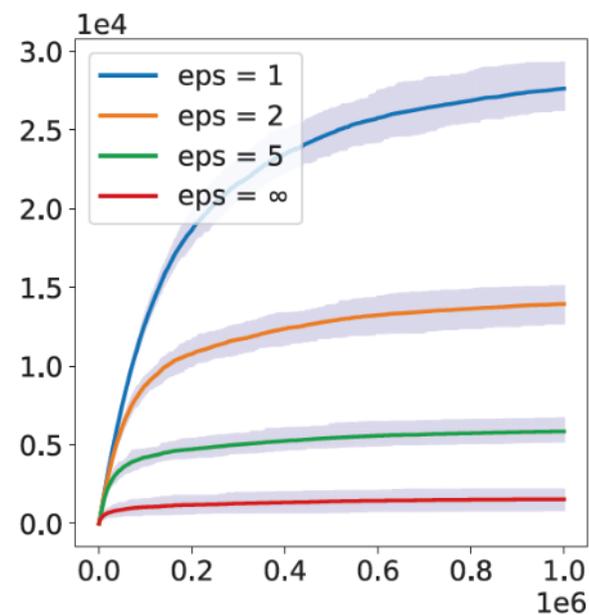
- Synthesize Fig. (a), Fig. (b)
- UCI Fig. (c)



(a)



(b)



(c)

Thanks for your attention!

ACKNOWLEDGEMENT

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